

# COMPILING PATTERN MATCHING TO IN-PLACE MODIFICATIONS



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# Our Starting Point

```
struct tree {  
    int a;  
    struct tree *l, *r;  
};
```

```
void mirror_left (struct tree *t) {  
    t && t->right = t->left;  
}
```

- + Performance
- + No copy
- + Flexibility
- Manual memory handling

```
type tree =  
| Empty  
| Node of tree * int * tree  
  
let mirror_left = function  
| Empty -> Empty  
| Node(l,o,r) -> Node(l,o,l)
```

- + Automatic memory handling
- Hard to finely tune
- No control on the layout of the type
- Copies

# Our Starting Point

```
struct tree {  
    int a;  
    struct tree *l, *r;  
};  
  
void  
}  
  
type tree =  
| Empty  
| Node of tree * int * tree
```

## How to take the best of both worlds?

- + No copy
- + Flexibility
- Manual memory handling
- Hard to finely tune
- No control on the layout of the type
- Copies

# A Language to Describe Structural Transformations



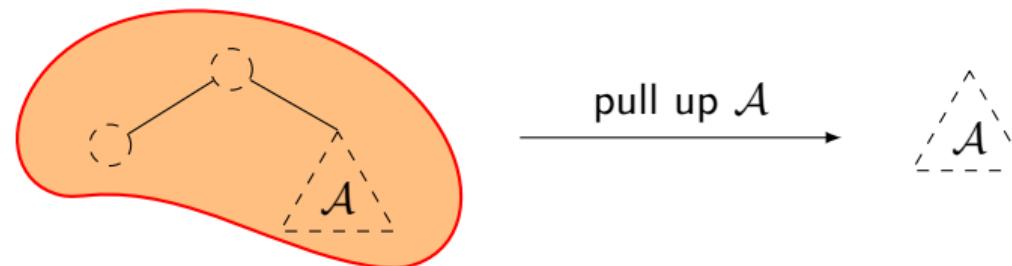
# A Language to Describe Structural Transformations



```
type tree = Empty | Node (tree,int,tree)
```

The tree is stored in an array

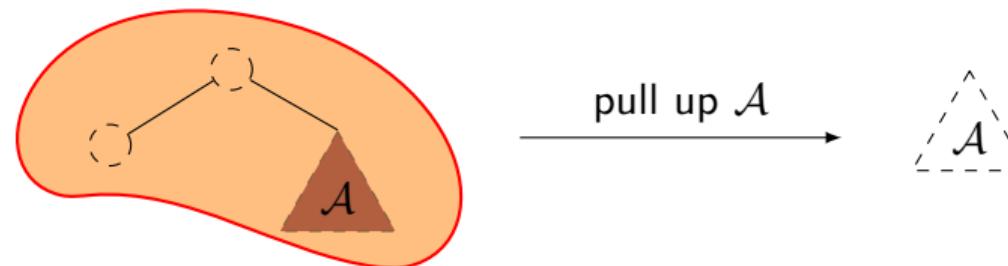
# A Language to Describe Structural Transformations



```
type tree = Empty | Node (tree,int,tree)
```

```
pull_up (t : tree) : tree = rewrite t {  
    | Node(a,i,Node(b,j,c)) -> Node(b,j,c)  
    | Empty -> Empty  
}
```

# A Language to Describe Structural Transformations



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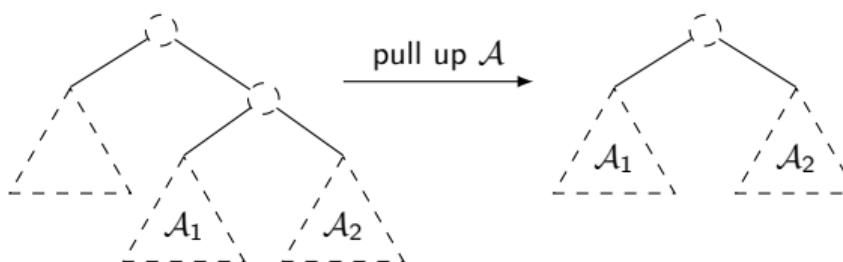


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# Step1: Compute Subtree Movements on Pull Up

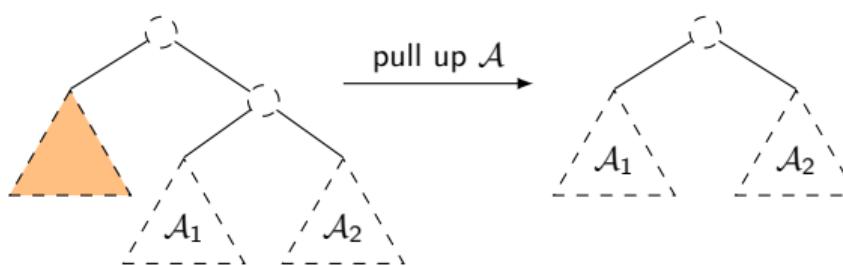
`Node( a , i ,Node( b , j , c )) -> Node( b , j , c )`



$(j : \text{int} | .2/\text{tree}.1/\text{int} \rightarrow .1/\text{int})$   
 $(i : \text{int} | .1/\text{int} \rightarrow \emptyset)$   
 $(b : \text{tree} | .2/\text{tree}.0/\text{tree} \rightarrow .0/\text{tree})$   
 $(a : \text{tree} | .0/\text{tree} \rightarrow \emptyset)$   
 $(c : \text{tree} | .2/\text{tree}.2/\text{tree} \rightarrow .2/\text{tree})$

# Step1: Compute Subtree Movements on Pull Up

$\text{Node}(a, i, \text{Node}(b, j, c)) \rightarrow \text{Node}(b, j, c)$



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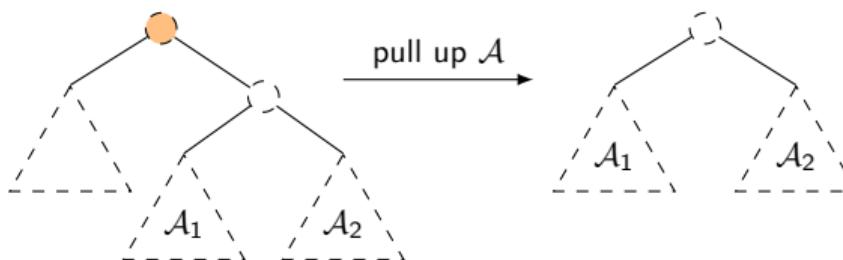
$(b : tree \mid .2/tree. 0/tree \rightarrow .0/tree)$

$(a : tree \mid .0/tree \rightarrow \emptyset)$

$(c : tree \mid .2/tree. 2/tree \rightarrow .2/tree)$

# Step1: Compute Subtree Movements on Pull Up

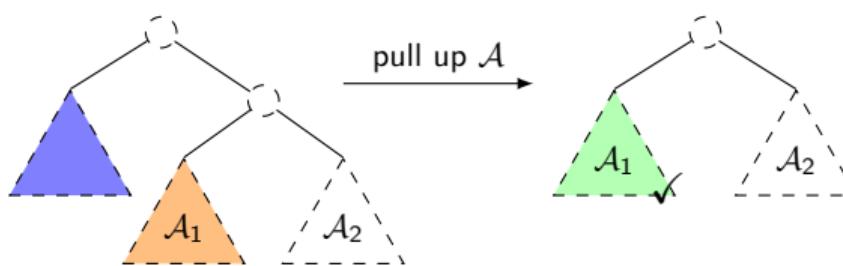
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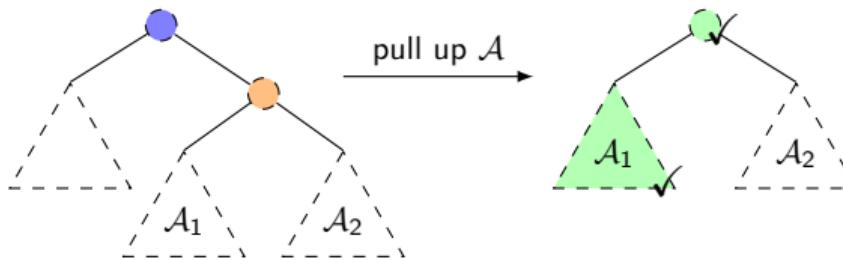
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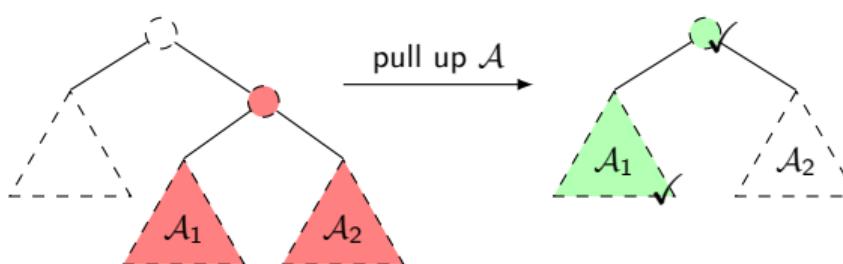
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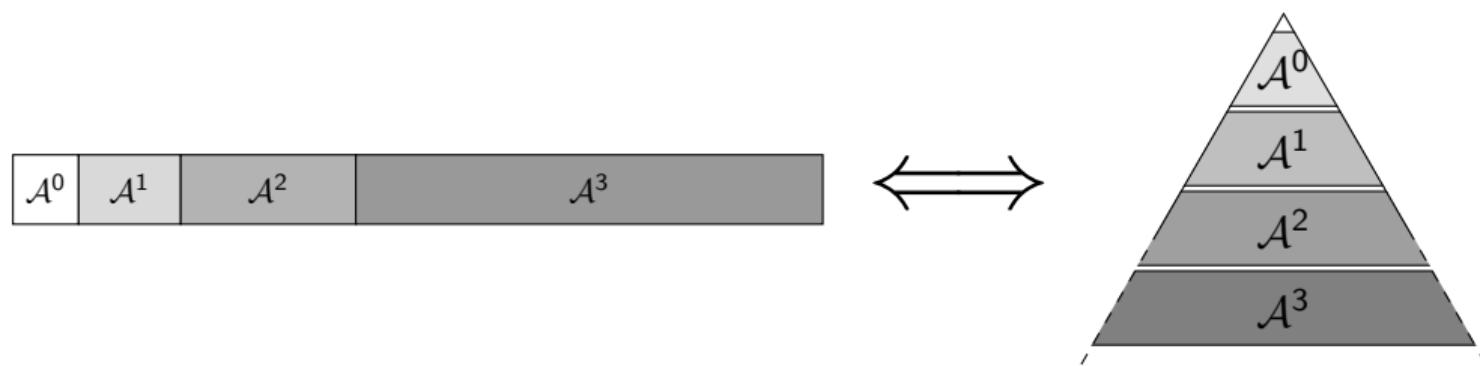
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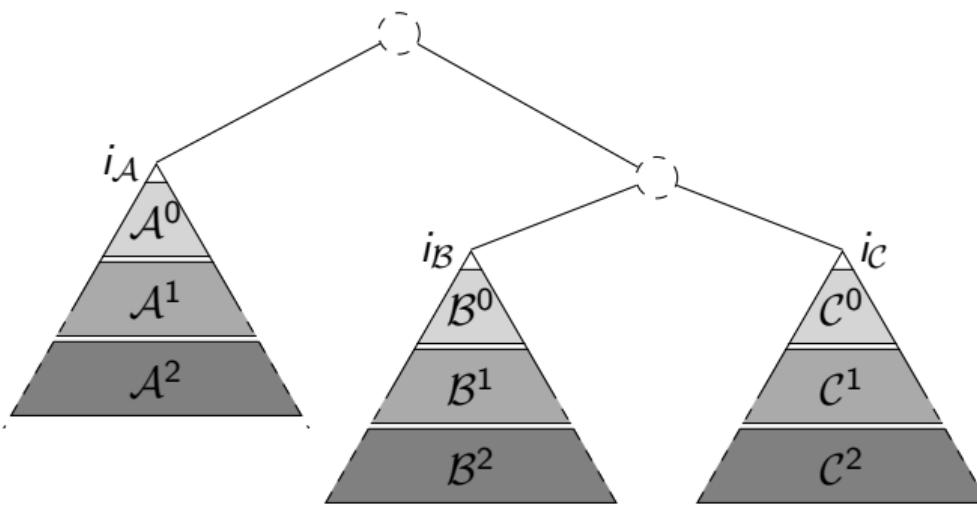
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# Use Layers to Subdivide Memory Movements on Pull Up



## Step 2: Layer-aware movements (1/2) on Pull Up

$\text{Node}(a, i, \text{Node}(b, j, c)) \rightarrow \text{Node}(b, j, c)$



$(. . 0/\text{tree. } \varphi^{k_0} \rightarrow \text{External})$  (a)

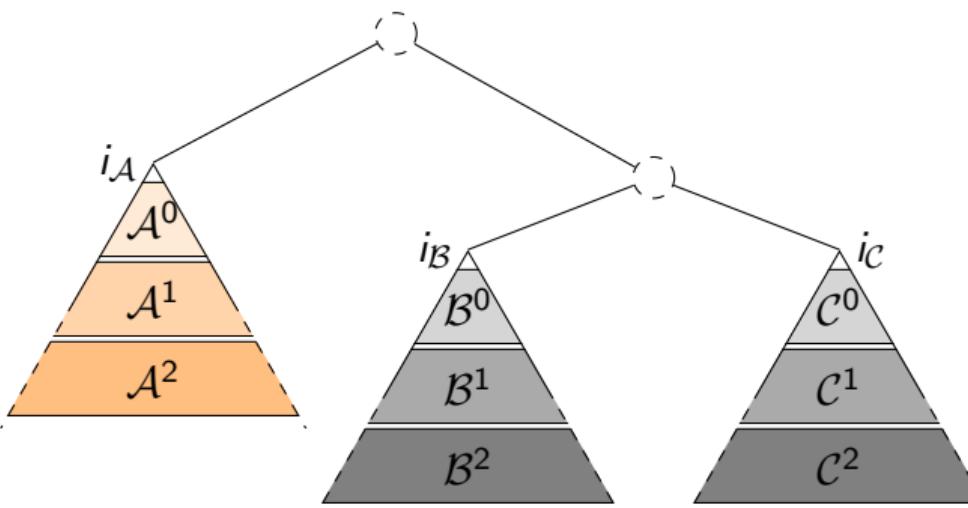
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$(. . 2/\text{tree. } 0/\text{tree. } \varphi^{k_1} \rightarrow . . 0/\text{tree. } \varphi^{k_1})$  (b)

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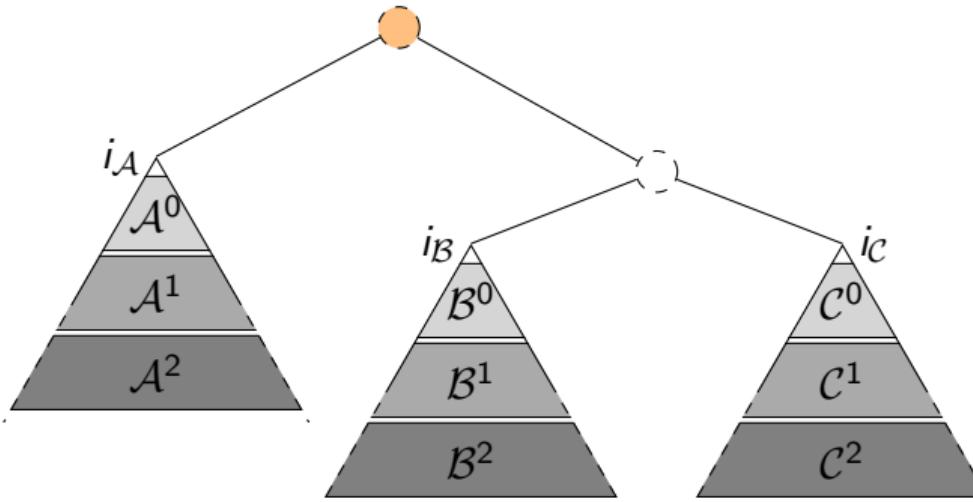
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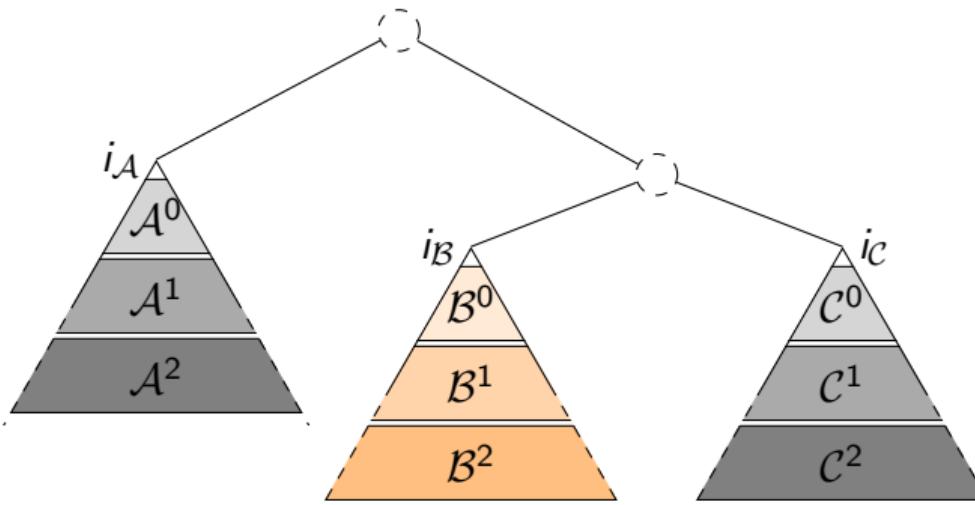
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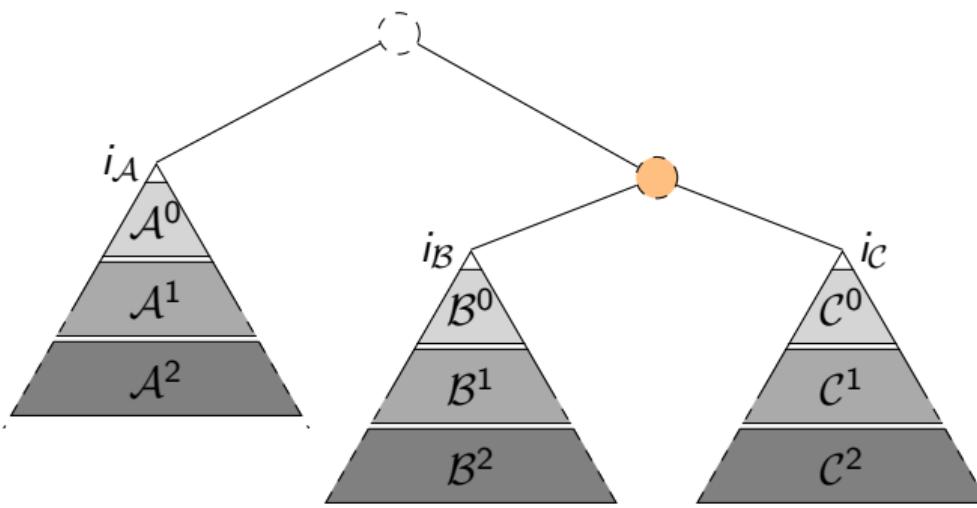
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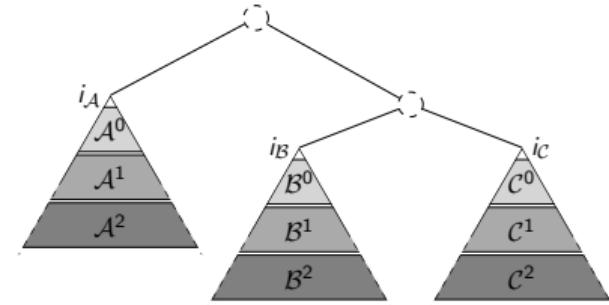
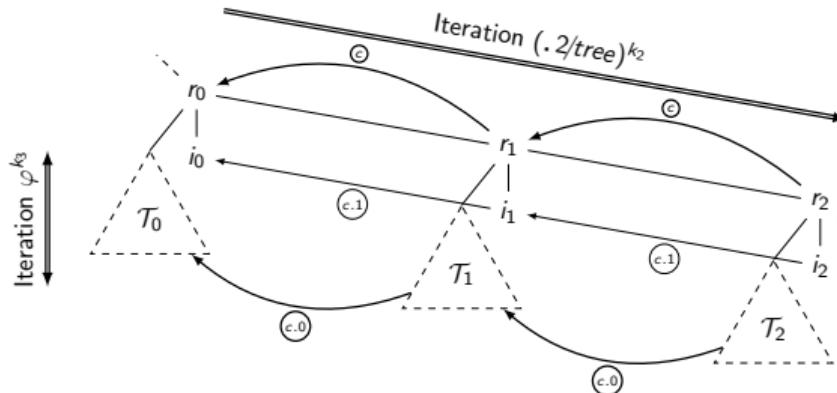
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## Step 2.5: Layer-aware movements (2/2) on Pull Up

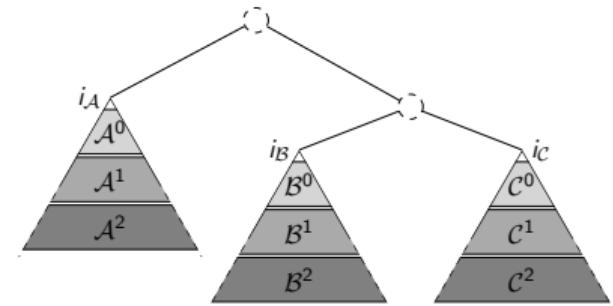
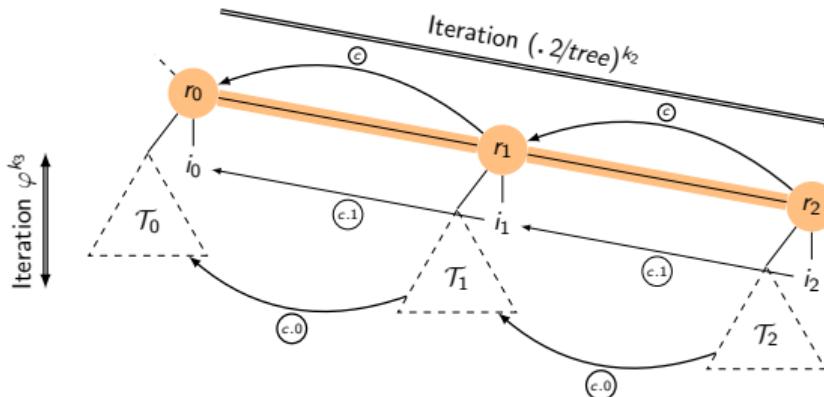


$$\langle\langle (.2/\text{tree})^{k_2+2} \rightarrow (.2/\text{tree})^{k_2+1} \rangle\rangle \quad (c)$$

$$\langle\langle (.2/\text{tree})^{k_2+2}.1/\text{int} \rightarrow (.2/\text{tree})^{k_2+1}.1/\text{int} \rangle\rangle \quad (c1)$$

$$\langle\langle (.2/\text{tree})^{k_2+2}.0/\text{tree}.\varphi^{k_3} \rightarrow (.2/\text{tree})^{k_2+1}.0/\text{tree}.\varphi^{k_3} \rangle\rangle \quad (c0)$$

## Step 2.5: Layer-aware movements (2/2) on Pull Up



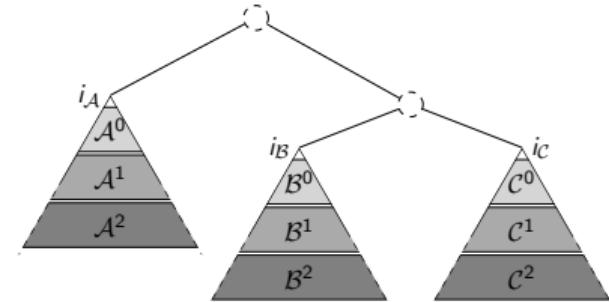
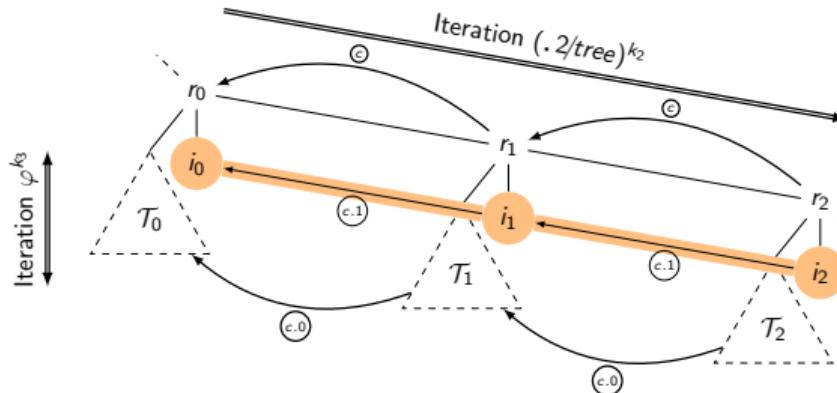
$$\langle \langle (.2/tree)^{k_2+2} \rightarrow (.2/tree)^{k_2+1} \rangle \rangle \quad (c)$$

$$\langle \langle (.2/tree)^{k_2+2} \cdot 1/int \rightarrow (.2/tree)^{k_2+1} \cdot 1/int \rangle \rangle \quad (c1)$$

$$\langle \langle (.2/tree)^{k_2+2} \cdot 0/tree \cdot \varphi^{k_3} \rightarrow (.2/tree)^{k_2+1} \cdot 0/tree \cdot \varphi^{k_3} \rangle \rangle \quad (c0)$$

## Fine-Grain Memory Movements

## Step 2.5: Layer-aware movements (2/2) on Pull Up

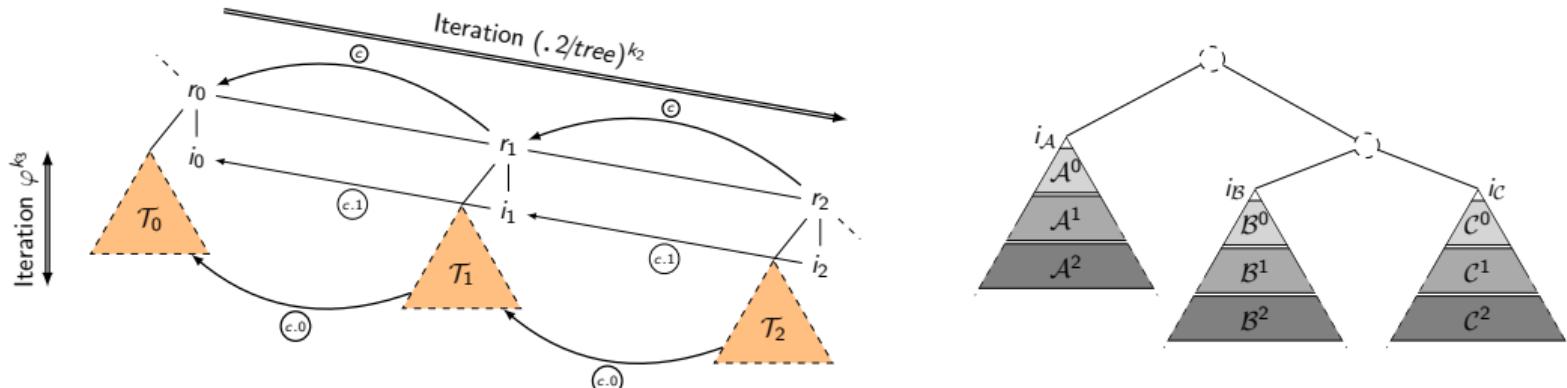


$$\langle\langle (.2/\text{tree})^{k_2+2} \rightarrow (.2/\text{tree})^{k_2+1} \rangle\rangle \quad (c)$$

$$\langle\langle (.2/\text{tree})^{k_2+2}.1/\text{int} \rightarrow (.2/\text{tree})^{k_2+1}.1/\text{int} \rangle\rangle \quad (c1)$$

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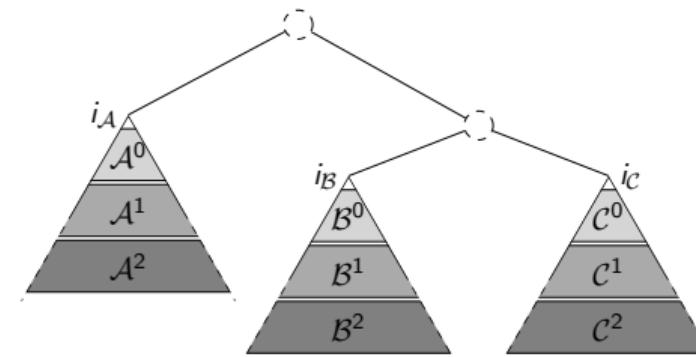
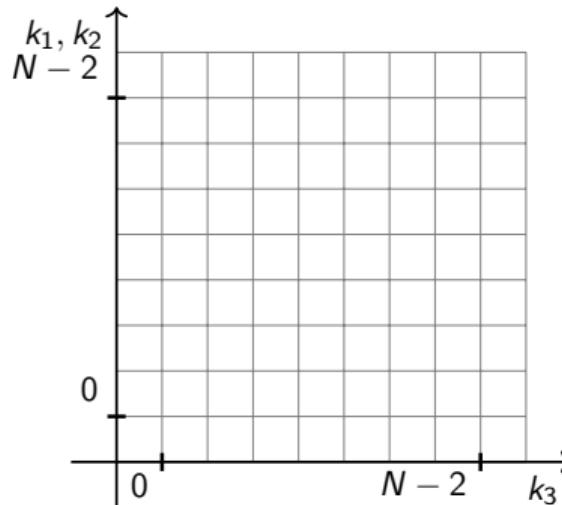
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$$\langle\langle (.2/tree)^{k_2+2}.1/int \rightarrow (.2/tree)^{k_2+1}.1/int \rangle\rangle \quad (c1)$$

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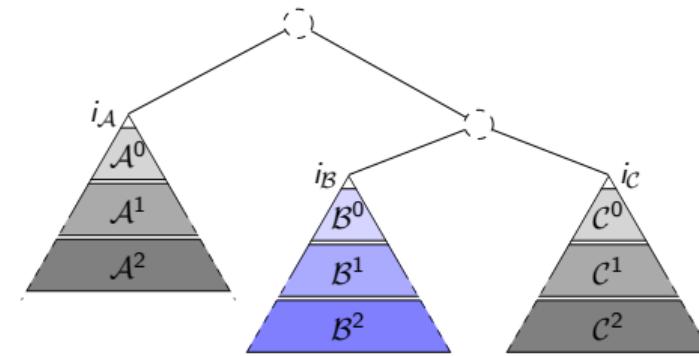
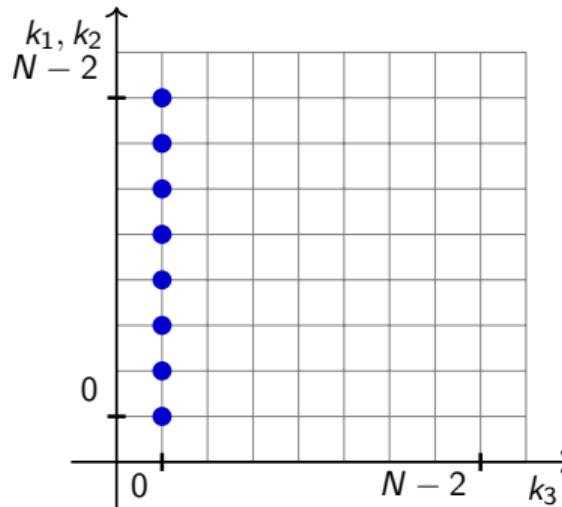
## Memory Movements Domains

## Step 3: Characterizing Memory Movements of Pull Up (Source)



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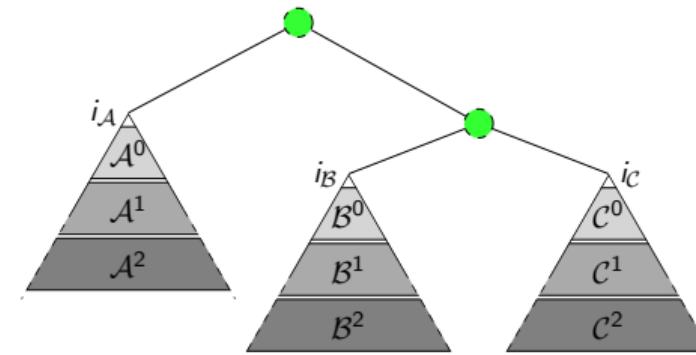
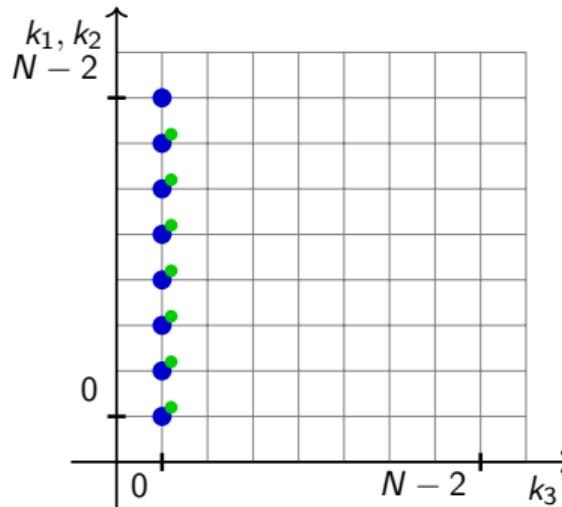


$$\langle .2/tree.0/tree.\varphi^{k_1} \rightarrow .0/tree.\varphi^{k_1} \rangle$$

(b)

## Memory Movements Domains

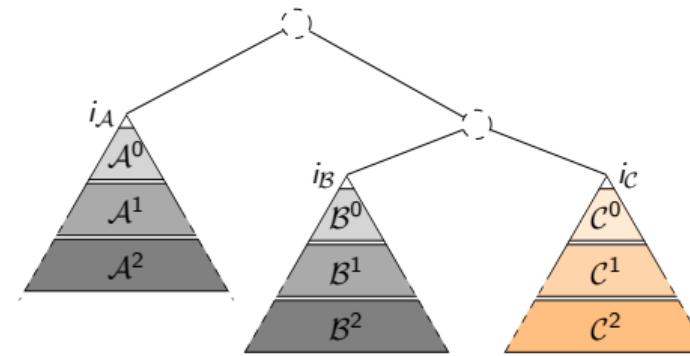
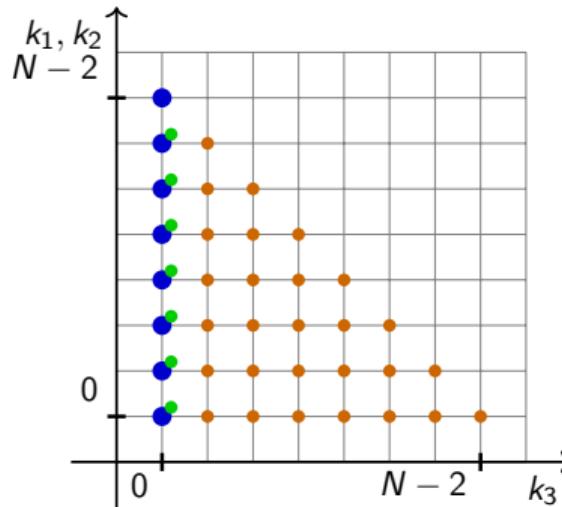
## Step 3: Characterizing Memory Movements of Pull Up (Source)



$$\langle (.2/tree)^{k_2+2} \cdot 1/int \rightarrow (.2/tree)^{k_2+1} \cdot 1/int \rangle \quad (c1)$$

## Memory Movements Domains

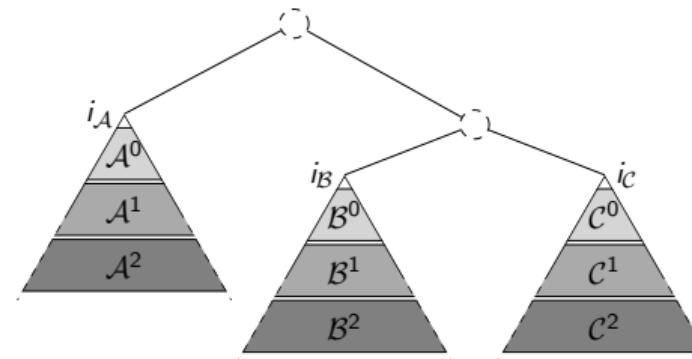
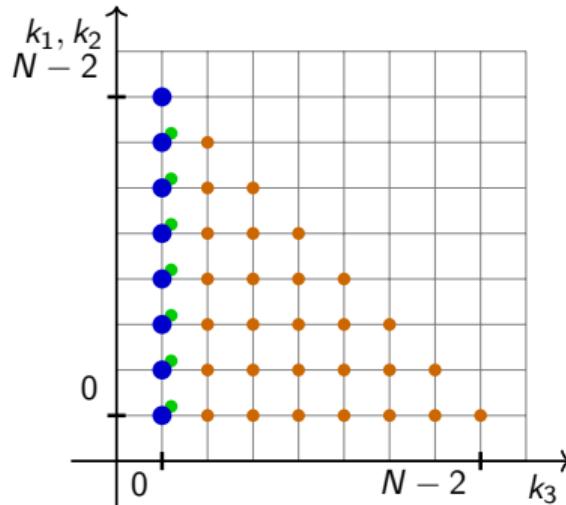
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$$\langle (.2/\text{tree})^{k_2+2}.0/\text{tree}.\varphi^{k_3} \rightarrow (.2/\text{tree})^{k_2+1}.0/\text{tree}.\varphi^{k_3} \rangle \quad (c0)$$

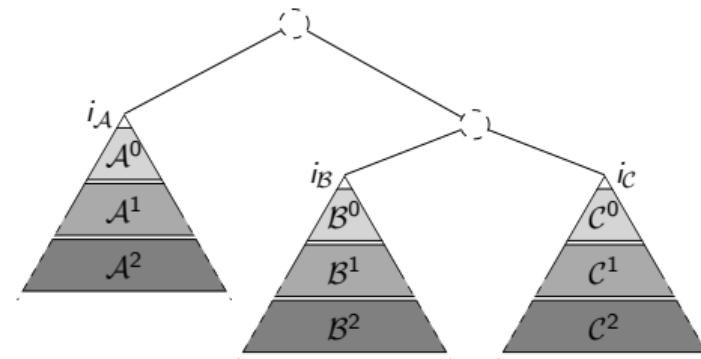
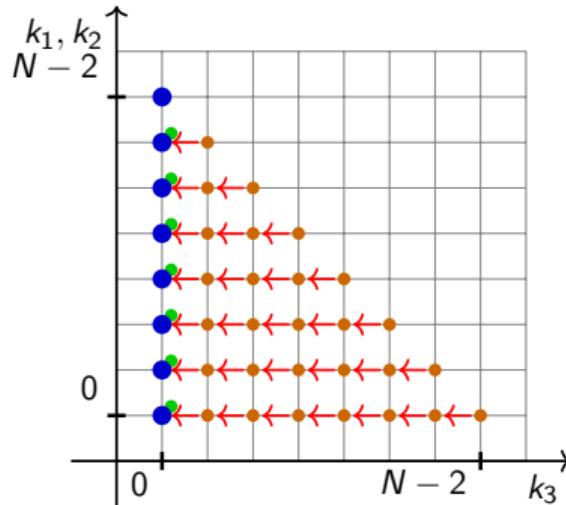
Inter-Movement Dependencies

## Characterizing Memory Movements of Pull Up (Dependencies)



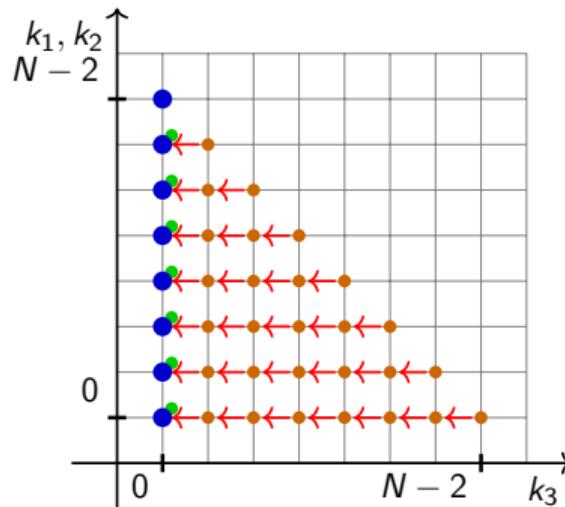
Inter-Movement Dependencies

## Characterizing Memory Movements of Pull Up (Dependencies)

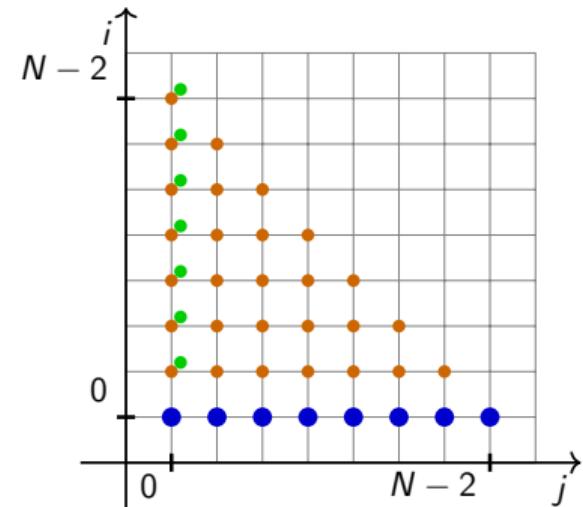


Schedule Computation

## Step 4: Scheduling as an Integer Linear Program Problem



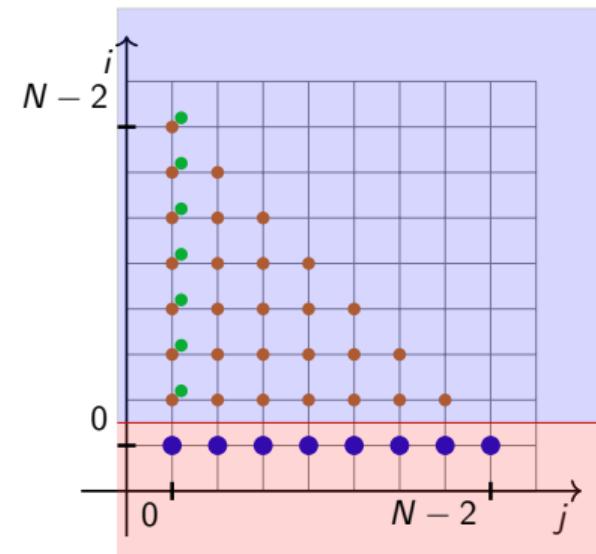
ILP Solver



Application of the Algorithm to the Running Example

## Step 5: From Schedule to Code

```
for (i = 0 ; i <= 0 ; i += 1) // P1
  for (j = 0 ; j <= N-2 ; j += 1)
    (.2/tree.0/tree.φj → .0/tree.φj) // b
for (i = 1 ; i <= N-2 ; i+= 1) // P2
  for (j = 0 ; j <= 0 ; j += 1) // P'1,2
    ((.2/tree)i+1.1/int → (.2/tree)i.1/int) // c1
for (j = 0 ; j <= N - i - 2 ; j += 1) // P'2,2
  ((.2/tree)i+1.0/tree.φj → (.2/tree)i.0/tree.φj) // c0
```



Application of Quilleré's Algorithm

# Conclusion & Future Work

- A promising technique to compile pattern matching to in-place transformation:
  - Based on term-rewriting
  - Use linear equations on regexps!
  - Schedule the operations (pipelining, parallelization)
  - Generate code

## Future Work

- Extend the input language
  - recursion
  - guards
- Improve the code generation
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# Contents

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- Structural Transformations
- Coarse Memory Movements
- Fine-Grain Memory Movements

## 2 Scheduling Memory Movements

- Memory Movements Domains
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- Schedule Computation

## 3 Code Generation

- Application of the Algorithm to the Running Example